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**Problems Solved:** 

Name:

Matrikel-Nr.:

Problem 36. Let

$$Y := (\lambda f.((\lambda x.(f(xx)))(\lambda x.(f(xx))))).$$

(just as in Section 3.2 of the Lecture Notes). Show by an explicit derivation that

 $(YF) \rightarrow^* (F(YF)).$ 

Problem 37. True or false?

- 1.  $(2n+3)(3n+2) = O(n^2)$
- 2.  $(2n+3) + \log_2(3n^6+2) = O(n)$
- 3.  $\frac{1024}{2^n} = O(1)$
- 4.  $\frac{1024}{2^n} = \Theta(\log_2(n))$
- 5.  $4^n = O(2^n)$
- 6.  $2^n = O(4^n)$

Prove your answers based on Definition 45 from the lecture notes.

**Problem 38.** Let  $f, g, h : \mathbb{N} \to \mathbb{R}_{\geq 0}$ . Prove or disprove based on Definition 45 from the lecture notes.

- 1. f(n) = O(f(n))
- 2.  $f(n) = O(g(n)) \implies g(n) = O(f(n))$
- 3.  $f(n) = O(g(n)) \land g(n) = O(h(n)) \implies f(n) = O(h(n))$

**Problem 39.** Write a LOOP program that computes the function  $f : \mathbb{N} \to \mathbb{N}$ ,  $f(n) = 2^n$ .

- 1. Count the number of variable assignments (depending on n) during the execution of your LOOP program with input n.
- 2. What is the time complexity of your program (depending on n)?
- 3. Is it possible to write a LOOP program with time complexity better than  $O(2^n)$ ? Give an informal reasoning of your answer.
- 4. Let l(k) denote the bit length of a number  $k \in \mathbb{N}$ . Let b = l(n), i.e., b denotes the bit length of the input. What is the time complexity of your program depending on b, if every variable assignment  $x_i := x_j + 1$  costs time  $O(l(x_j))$ ?

## Berechenbarkeit und Komplexität, WS2015

**Problem 40.** Let  $\Sigma = \{0,1\}$  and let  $L \subseteq \Sigma^*$  be the set of binary numbers divisible by 3, i.e.,

$$L = \{x_n \dots x_1 x_0 : 3 \text{ divides } \sum_{k=0}^n x_k 2^k\}.$$

(By convention, the empty string  $\varepsilon$  denotes the number 0 and so it is in L too.)

- 1. Design a Turing machine M with input alphabet  $\Sigma$  which recognizes L, halts on every input, and has (worst-case) time complexity T(n) = n. Write down your machine formally. (A picture is not needed.) *Hint:* Three states  $q_0, q_1, q_2$  suffice. The machine is in state  $q_r$  if the bits read so far yield a binary number which leaves a remainder of r upon division by 3. The transition from one state to another represents a multiplication by 2 and the addition of 0 or 1.
- 2. Determine S(n),  $\overline{T}(n)$  and  $\overline{S}(n)$  for your Turing machine.
- 3. Is there some faster Turing machine that achieves  $\overline{T}(n) < n$ ? (Justify your answer.)