

# Formal Methods in Software Development

## Sample Exam

Wolfgang Schreiner  
Wolfgang.Schreiner@risc.jku.at

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**Last Name:**

**First Name:**

**Matrikelnummer:**

**Studienkennzahl:**

KV3 (questions 1-4): 85 points total.

KV4 (questions 1-5): 100 points total

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**Please indicate whether you want to be graded for KV3 or KV4.**

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1. (20 points) Write a heavy-weight JML specification for the following method of the Java library (the specification shall be as expressive as possible).

```
public static void fill(int[] a, int fromIndex, int toIndex, int val)
```

Assigns the specified int value to each element of the specified range of the specified array of ints. The range to be filled extends from index fromIndex, inclusive, to index toIndex, exclusive. (If fromIndex==toIndex, the range to be filled is empty.)

Parameters:

a - the array to be filled

fromIndex - the index of the first element (inclusive) to be filled with the specified value

toIndex - the index of the last element (exclusive) to be filled with the specified value

val - the value to be stored in all elements of the array

Throws:

IllegalArgumentException - if fromIndex > toIndex

ArrayIndexOutOfBoundsException - if fromIndex < 0 or toIndex > a.length



2. (20 points) Derive the strongest postcondition of the command  $c$

```
if (i < 10)
{
  a[i] = a[i]+3;
  i = i+1;
}
```

for precondition  $a[2] = 5$  (ignoring ‘index out of bound’ violations).

Then derive a judgement of form  $c : [F]^{x,\dots}$  for some state transition  $F$  and variable frame  $\{x, \dots\}$ .

In both cases, show all derivation steps and finally simplify the derived formulas as far as possible.

3. (20 points) Take the following program which is supposed to compute for given  $n \in \mathbb{N}$  the result  $s := n^2$ :

```
{n = oldn}
  s = 0; i = 1;
  while (i <= n)
  {
    s = s+2*i-1;
    i = i+1;
  }
{s = n2 ∧ n = oldn}
```

- Assume you are given a suitable loop invariant  $I$  and termination term  $T$ ; using  $I$  and  $T$  state all verification conditions (classical logic formulas) that have to be proved for verifying partial correctness and termination of the program.
- Construct for input  $n = 5$  a table with the values of the variables before/after each loop iteration. Using this table as a hint, give suitable definitions for  $I$  and  $T$  and perform the verification.



4. (25 points) Take the following asynchronous composition of two processes operating on shared variables  $x, y, i, j$ :

```
initially x = y = i = 0 and j = 1
loop      ||      loop
  x = x+j;  ||      wait i > 0;
  i = 1-i;  ||      y = y+i;
           ||      j = 1-j;
```

- (10 points) Give a formal model of the system (using the interleaving assumption for asynchronous composition); do not forget to model the program counters of the two processes.
- (5 points) Formalize in LTL the property “if at any time  $i$  becomes greater than zero, then eventually also  $y$  will become greater than zero”.
- (10 points) Is this property true for above system? If yes, explain why. If not, show an execution trace that violates the property.

In the second case, if there exists a fairness assumption for the system execution, under which the property holds, state this assumption and explain in detail why it makes the property true.





5. (15 Points)

Derive the performance measures of an M/M/1 system (i.e., justify your answers).