Formal Methods in Software Development
Assignment 2 (November 16)

Wolfgang Schreiner
Wolfgang.Schreiner@risc.jku.at

The result is to be submitted by the deadline stated above *via the Moodle interface* of the course as a `.zip` or `.tgz` file which contains

1. a PDF file with
   - a cover page with the course title, your name, Matrikelnummer, and email address,
   - a section for each part of the exercise with the requested deliverables and optionally any explanations or comments you would like to make;
2. the RISCAL specification (`.txt`) file(s) used in the exercise.

Email submissions are *not* accepted.

Assignment 2: Deriving and Checking Verification Conditions

Given an array a of non-negative integers and a positive integer n, find a position of the maximum of the first n elements of a, i.e., a non-negative integer p less than n such that the element of a at p is that maximum. We claim that this specification is implemented by the following program fragment:

```
p := -1; e := 0; i := 0;
while i < n do
{
  if a[i] ≥ e then { p := i; e := a[p]; }
  i := i+1;
}
```

The goal of this assignment is to verify this claim by deriving and checking the verification conditions whose validity implies the total correctness of this code with respect to the given specification.

1. Take file `maximum.txt` which embeds above code in a procedure `maximumPosition` and equip this procedure with suitable preconditions (`requires`) and postconditions (`ensures`) that formalize above specification. Check (for moderately large $N \geq 4$ and $M \geq 3$) that the procedure satisfies the specification.

   Hint: in order to avoid any later confusion between the program variable i and mathematical variables, it is recommended to name quantified variables different from i.

2. Select the operation button "Show/Hide Tasks" to display all tasks related to the specification of the procedure ("Execute specification", "Validate specification" and "Verify specification preconditions"). Validate the specification by executing these tasks (run "Execute specification" with execution option "Silent" switched off to investigate the input/output pairs allowed by your specification; the checks of the other tasks which denote theorems may be performed with option "Silent" switched on). If a task stays red, i.e., the corresponding theorem does not hold, you may choose the entry "Show Counterexample" from the popup menu of the task (which appears with a right-click) to get some further insight, which may help you to improve your procedure specification.

3. Annotate the loop with suitable invariants (`invariant`) and termination term (`decreases`). Rerun the procedure check to ensure that your annotations are not too strong (but they may be still too weak to carry the verification).

   Please note that RISCAL treats procedure parameters such as a, m and n as unmodifiable constants; thus it is not necessary to specify that their values remain unchanged. Furthermore, all the knowledge from the precondition of the procedure is automatically inherited and does not have to specified in the invariants again. However, you have to express in the invariant the following information:

   - The range of i: $(\ldots \leq i \leq \ldots)$.
   - The range of p (please note the special case $i = 0$): (if $i = 0$ then $\ldots$ else $\ldots$).

- The value of e (again please note the special case $i = 0$).
- The knowledge about p that arises from the postcondition (which talks about the situation after the termination of the loop when the whole array has been processed) and is adequately generalized (to include the situation before/after every loop iteration when only part of the array has been processed): $(\forall k : 0 \leq k < i \Rightarrow \ldots)$.

4. Now demonstrate your knowledge of the Hoare calculus by *manually* deriving from the specification and the loop annotations the verification conditions whose validity implies the total correctness (partial correctness *and termination*) of the program.

   Hint: Hoare calculus reasoning yields five conditions: one for showing that the input condition of the loop (which is different from the input condition of the program!) implies the invariant, one for showing that the invariant and the negation of the loop condition implies the output condition, two for showing that the invariant is preserved and the value of the termination term is decreased for each of the two possible execution paths in the loop body, one for showing that the invariant implies that the value of the termination term does not become negative (if you apply weakest precondition reasoning within the loop body, only one condition is derived from the loop body).

   Do not only give the final verification conditions but show in detail their derivation by application of Hoare calculus (respectively the predicate transformer calculus). *Don't try to "guess" the conditions!*

   Check these conditions with RISCAL, in the style of the verification of the "linear search" algorithm presented in class. For this purpose, define predicates `Input`, `Output`, and `Invariant` and a function `Termination`, where (as shown in class) `Invariant` and `Termination` should be parametrized over the program variables. Then define five theorems `A`, `T`, `B1`, `B2`, `C` describing the verification conditions and check these. Do not forget to make the preconditions of the procedure also preconditions of these theorems. If a theorem does not hold, you may select the operation button "Show/Hide Tasks" and choose the menu entry "Show Counterexample" to get some further insight, which may help you to correct your annotations.

   Also attempt the theorem proving capabilities of RISCAL (menu option "Apply Theorem Prover") to verify the validity of these conditions for *arbitrary* values of M and N with proof method SMT (make sure that in menu "TP" menu entry "Method SMT" is selected).

5. Then apply the capabilities of RISCAL to automatically generate the necessary verification conditions from the annotated program. For this select the operation button "Show/Hide Tasks" and check all tasks related to the implementation ("Verify specification preconditions", "Verify implementation preconditions", "Verify correctness of result", "Verify iteration and recursion"). If your annotations are adequate (strong enough but not too strong), then all red tasks turn blue (as demonstrated in class for the "linear search" example). Again, if a task stays red, i.e., the corresponding verification condition does not hold, you may choose the entry "Show Counterexample" to get some further insight, which may help you to correct your annotations.

Again attempt the theorem proving capabilities of RISCAL to verify the validity of the verification conditions for *arbitrary* values of M and N with proof method SMT.

The deliverables for this assignment consists of the following items:

1. a nicely formatted copy of the RISCAL specification (included as text, not as screenshots);
2. a detailed manual derivation of the verification conditions;
3. for each check of a verification condition a reasonable selection of the output (included as text, not as screenshots) with an explicit statement whether the check has succeeded; if a check failed, give a conjecture why this is the case.
4. screenshots of (part of) the RISCAL user interface that show all automatically generated tasks after checking (panel "Tasks", all/most of these tasks should be blue);
5. an explicit statement of whether all tasks could be successfully checked or not for some fixed values of M and N; if some task could not be successfully checked, show its definition (menu entry "Print definition") and your conjecture of the reason of the failure;
6. an explicit statement of whether all theorems could be successfully proved for arbitrary values of M and N; if some theorem could not be successfully proved, show (a reasonable selection) of the output of the menu entry "Print Prover Output" and your conjecture of the reason of the failure.