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Computational Logic, WS 2024/2025,  
Exercise sheet 9,  
due date: **19 January 2025, 23:59 via Moodle**

**Problem 39 (30 Points)**

Let  $eq$  be a binary predicate such that  $(\forall x eq(x, x))$ ,  $(\forall x \forall y eq(x, y) \Rightarrow eq(y, x))$  and  $(\forall x \forall y \forall z eq(x, y) \wedge eq(y, z) \Rightarrow eq(x, z))$ . Consider the following formulas:

- a)  $(\exists x \exists y \exists z \forall t eq(t, x) \vee eq(t, y) \vee eq(t, z))$
- b)  $(\forall x \exists y \exists z \exists t \neg eq(x, y) \wedge \neg eq(x, z) \wedge \neg eq(x, t) \wedge \neg eq(y, z) \wedge \neg eq(y, t) \wedge \neg eq(z, t))$

Using the tableau procedure with free variables and unification prove that the conjunction of the two formulas is unsatisfiable.

**Problem 40 (20 Points)**

Let  $x, y, z$  and  $t$  be variables;  $a$  and  $b$  constants. Consider the following set of clauses:

$$\{\neg P(f(x, y), z) \vee R(z), Q(x) \vee \neg R(a), P(t, a) \vee R(a), \neg Q(b)\}$$

- a) Prove by first-order resolution that this set is unsatisfiable.
- b) Translate this set of clauses into both a set of Horn clauses (using the concrete syntax of facts and rules) and into a Prolog program, if possible; otherwise, explain why such translations are not possible.

**Problem 41 (20 Points)**

Consider the following logic program:

$r(g(X)) \leftarrow t(X, Y, f(X))$

$t(a, b, f(a)) \leftarrow$

$p(V, W) \leftarrow r(V)$

with query  $\leftarrow p(U, b)$ .

What will happen when using SLD resolution?

**Problem 42 (30 Points)**

Formalize each group of sentences, then prove that the last one is a consequence of the others using resolution and paramodulation.

- a)
  - Rex is a dog whose owner lives in a white house.
  - Rob is the only person who lives in a white house.
  - Therefore, Rob is the owner of Rex.
- b)
  - Rob owns only white dogs.
  - Rex and Daisy have the same owner.
  - Rob is the owner of Daisy.
  - Therefore, Rex is a white dog.