## Formal Specification of Abstract Datatypes Exercise 5 (June 27)

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The result is to be submitted by the deadline stated above via the Moodle interface as a .zip or .tgz file which contains

- a PDF file with
  - a cover page with the title of the course, your name, Matrikelnummer, and emailaddress,
  - the formal specifications in the style of "Thinking Programs",
  - the CafeOBJ specifications,
  - comprehensive tests of the CafeOBJ specifications (sample reductions),
  - optionally any explanations or comments you would like to make;
- the CafeOBJ (.mod) file(s) of the specifications.

## **Exercise: Directed Graphs**

In set theory, a *directed graph* is a pair G = (N, E) where N is a set of elements called *nodes* and E is a set of pairs of nodes called *edges*.

The goal of this exercise is to specify an abstract datatype of directed graphs with associated operations where  $N = \mathbb{N}_n = \{0, ..., n - 1\}$  for some  $n \in \mathbb{N}$ .

**Specification** First develop a formal specification of the following shape:

```
spec GRAPH import NAT :=
... {
  type Graph = g(Nat,Edges)
   type Edges = ...
}
then ... {
   graph: Nat \rightarrow Graph
   link: Graph \times Node \times Node \rightarrow Graph
  nodes: Graph \rightarrow Nat
   edges: Graph \rightarrow Nat
   adjacent \subseteq Graph \times Node \times Node
   connected \subseteq Graph \times Node \times Node
   distance: Graph \times Node \times Node \rightarrow Nat
   complete \subseteq Graph
   connected \subseteq Graph
   cyclic \subseteq Graph
   . . .
}
```

Make sure that two graphs are equal if they have the same nodes and the same edges.

Here graph(n) denotes the graph with *n* nodes and no edges while  $link(g, n_1, n_2)$  denotes the graph that adds to graph *g* an edge from  $n_1$  to  $n_2$  (if  $n_1$  or  $n_2$  is not a node of *g* or *g* already contains that edge, the graph remains unchanged). Furthermore, nodes(g) denotes the number of nodes in *g* while edges(g) denotes the number of its edges.

The predicate  $adjacent(g, n_1, n_2)$  holds if g has an edge from  $n_1$  to  $n_2$  (thus it does trivially not hold, if  $n_1$  or  $n_2$  is not a node of g). The predicate  $connected(g, n_1, n_2)$  holds if there exists for some  $k \ge 1$  a path of nodes  $n_1 = m_1, m_2, \ldots, m_k = n_2$  where every pair  $m_i$  and  $m_{i+1}$  with i < k is adjacent (thus every node n is trivially connected to itself by the path  $n = m_1 = n$ ). For two connected nodes  $n_1$  and  $n_2$  in g,  $distance(g, n_1, n_2)$  denotes the number of edges in the shortest path from  $n_1$  to  $n_2$  (thus every node has distance 0 to itself).

The predicate complete(g) holds if every pair of nodes in g is adjacent. The predicate connected(g) holds if every pair of nodes in g is connected. The predicate cyclic(g) holds, if there is some node in g that is connected to itself by a path with at least one edge.

Does your specification give rise to a monomorphic datatype (justify your answer)?

**CafeOBJ** Implement your specification by a tight CafeOBJ module

```
module! GRAPH
{
   protecting (NAT)
   ...
}
```

Test the CafeOBJ module with several reductions. Give the input and output of each test and your interpretation of the results (do they indicate errors in your specification or not?). If your specification contains errors, use the trace facilities of CafeOBJ for debugging.

Does your CafeOBJ module denote the same datatype as the one of the previous specification? Justify your answer.