## Problems Solved:

| 11 | 12 | 13 | 14 | 15 |
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## Name:

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Problem 11. Let $M_{1}=\left(Q_{1}, \Sigma, \delta_{1}, q_{1}, F_{1}\right)$ and $M_{2}=\left(Q_{2}, \Sigma, \delta_{2}, q_{2}, F_{2}\right)$ be two DFSM over the alphabet $\Sigma$. Let $L\left(M_{1}\right)$ and $L\left(M_{2}\right)$ be the languages accepted by $M_{1}$ and $M_{2}$, respectively.
Construct a DFSM $M=(Q, \Sigma, \delta, q, F)$ whose language $L(M)$ is the intersection of $L\left(M_{1}\right)$ and $L\left(M_{2}\right)$. Write down $Q, \delta, q$, and $F$ explicitly.
Hint: $M$ simulates the parallel execution of $M_{1}$ and $M_{2}$. For that to work, $M$ "remembers" in its state the state of $M_{1}$ as well as the state of $M_{2}$. This can be achieved by defining $Q=Q_{1} \times Q_{2}$.
Demonstrate your construction with the following DFSMs.


Problem 12. Let $m>1$ be a natural number and $r=\left(a^{*} b\right)^{m}$ be a regular expression.
Definition: If $a$ is a regular expression and $n$ is a natural number, then $a^{n}=$ $a \cdots a$ is the regular expression that results from the $n$-fold concatenation of $a$. For example: $a^{3}=a a a$.
Let $L_{1}$ and $L_{2}$ be the languages defined as follows: $L_{1}:=L(r), L_{2}:=\left\{\left(a^{k} b\right)^{m} \mid k \in \mathbb{N}\right\}$. Are $L_{1}$ and $L_{2}$ regular languages? Provide solid arguments to your answers.

Problem 13. Let $M_{1}$ be the DFSM with states $\left\{q_{0}, q_{1}, q_{2}\right\}$ whose transition graph is given below. Determine a regular expression $r$ such that $L(r)=L\left(M_{1}\right)$. Show the derivation of the the final result by the technique based on Arden's Lemma (see lecture notes).


Problem 14. Let $r$ be the following regular expression.

$$
a \cdot a \cdot(b \cdot a)^{*} \cdot b \cdot b^{*}
$$

Construct a nondeterministic finite state machine $N$ such that $L(N)=L(r)$. Show the derivation of the result by following the technique presented in the proof of the theorem Equivalence of Regular Expressions and Automata (see lecture notes).

Problem 15. Show that the language $L=\left\{a^{m} b^{n} \mid m, n \in \mathbb{N} \wedge m \geq 2 n\right\}$ is not regular.

