**Problems Solved:** 

 $31 \ 32 \ 33 \ 34 \ 35$ 

Name:

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**Problem 31.** Let a language  $L = L(T) \subseteq \{0,1\}^*$  be given by the code of a Turing machine  $\langle T \rangle$ . It is known that  $\varepsilon \in L$ .

Let  $S_0$  be the set of Turing machines of the form  $(Q, \{0, 1, X, \sqcup\}, \sqcup, \{0, 1\}, q_0, \emptyset)$ . Let  $S_1$  be the set of Turing machines of the form  $(Q, \{0, 1, X, \sqcup\}, \sqcup, \{0, 1\}, q_0, Q)$ . Is it decidable whether L = L(M) and  $M \in S_0$ ? That is: Is there a Turing machine  $D_0$  such that it takes a word w as input and returns "yes" if  $w = \langle M \rangle$  for a TM  $M \in S_0$  with the property L(M) = L, and returns "no" otherwise? What is the answer, if you replace  $S_0$  by  $S_1$ ? Justify your answers.

**Problem 32.** Let  $\Sigma$  be an alphabet and A be a language over  $\Sigma$  ( $A \subseteq \Sigma^*$ ). Let also A be semi-decidable, but not decidable. Prove that the complement of A, i. e.,  $\overline{A} = \Sigma^* \setminus A$ , is not decidable.

**Problem 33.** Let *L* be a finite language over an alphabet  $\{0, 1\}$ . Is the following problem (with input  $\langle M \rangle$ )

For a Turing maschine M it holds  $L(M) \supseteq L$ .

in general semi-decidable? Is it also in general decidable? Justify your answers.

**Problem 34.** Which of the following problems are decidable? In each problem below, the input of the problem is the code  $\langle M \rangle$  of a Turing machine M with input alphabet  $\{0, 1\}$ .

- (a) Does M have at least 4 states?
- (b) Is  $L(M) \subseteq \{0, 1\}^*$ ?
- (c) Is L(M) recursive?
- (d) Is L(M) finite?
- (e) Is  $10101 \in L(M)$ ?
- (f) Is L(M) not recursively enumerable?
- (g) Does there exist a word  $w \in L(M)$  such that M does not halt on w?

Justify your answer.

**Problem 35.** Show that the Acceptance Problem is reducible to the restricted Halting problem. First explain clearly which Turing machine you have to construct to prove this statement and then give a reasonably detailed description of this construction.

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