**Problems Solved:** 

16 | 17 | 18 | 19 | 20

Name:

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**Problem 16.** Construct a Turing machine  $M = (Q, \Gamma, \sqcup, \{0, 1\}, \delta, q_0, F)$  such that  $L(M) = \{1^k 0 1^{k+1} | k \in \mathbb{N}\}$ . Write down  $Q, \Gamma, F$  and  $\delta$  explicitly.

**Problem 17.** Write a RAM program that from a given natural number n prints its binary representation. In order to simplify the problem the output shall be in low positions first format, i.e., the number  $8_{10}$  is  $0001_2$  but not  $1000_2$ . Hint: please note that the computation of the quotient respectively remainder of a division by 2 can be implemented by the repeated subtraction of 2.

## Problem 18.

1. Show by using *only* the Definition of a *loop program* (Def. 23 in the lecture notes, Section 3.2.2) that the function

$$s(x_1, x_2) = \begin{cases} 1 & \text{if } x_1 < x_2, \\ 0 & \text{otherwise} \end{cases}$$

is loop computable. I.e. give an explicit loop program for s.

Note that it is not allowed to use an abbreviation like

 $x\,i\ :=\ x\,j\ -\ xk\,;$ 

2. Write a loop program that computes the function  $d : \mathbb{N} \to \mathbb{N}$  where  $d(x_1, x_2)$  is  $k \in \mathbb{N}$  such that  $k \cdot (x_2 + 1) = x_1 + 1$  if such a k exists. The result is  $d(x_1, x_2) = 0$ , if a k with the above property does not exist. For simplicity in the program for d, you are allowed to use a construction

For simplicity in the program for u, you are anowed to use a construction like the following (with the obvious semantics) where P is an arbitrary loop program.

IF xi < xj THEN P END;

**Problem 19.** Provide a loop program that computes the function  $f(n) = \sum_{k=1}^{n} k(k+1)$ , and thus show that f is loop computable. You are only allowed to use the constructs given in Definition 23 of the lecture

You are only allowed to use the constructs given in Definition 23 of the lecture notes.

**Problem 20.** Suppose P is a while-program that does not contain any WHILE statements, but might modify the values of the variables  $x_1$  and  $x_2$ .

Transform the following program into Kleene's normal form.

*Hint:* first translate the program into a goto program, replace the GOTOs by assignments to a control variable, and add the WHILE wrapper.

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 $\begin{array}{rl} & P\,;\\ & \textbf{END};\\ \textbf{END};\\ x0 \ := \ x0 \ + \ 1 \end{array}$ 

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