Problems Solved:

21 | 22 | 23 | 24 | 25

Name:

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Problem 21. We know that the function $p : \mathbb{N}^3 \to \mathbb{N}$ defined by $p(a, b, n) = a \uparrow^n b$ is not primitive recursive. However, the function $p_2 : \mathbb{N}^2 \to \mathbb{N}$, defined by $p_2(a, b) = a \uparrow^2 b$ is primitive recursive.

Show that fact by defining p_2 explicitly from the base functions, the (primitive recursive) function $\varepsilon(x, y) = x^y$, composition, and the primitive recursion scheme.

Problem 22. Let $f : \mathbb{N} \to \mathbb{N}$ be the (partial) function

$$f(x) = \begin{cases} y & \text{such that } x = y^2 \text{ if such a } y \text{ exists,} \\ \text{undefined} & \text{if there is no } y \text{ with } x = y^2. \end{cases}$$

- 1. Is f loop computable? (Justify your answer.)
- 2. Is f a primitive recursive function? (Justify your answer.)
- 3. Define f by using the base functions, composition, the primitive recursion scheme, and μ -recursion. Additionally you are allowed to use the (primitive recursive) functions

$$m: \mathbb{N}^2 \to \mathbb{N}, \quad (x, y) \mapsto x \cdot y$$

and $u: \mathbb{N}^2 \to \mathbb{N}$,

$$u(x,y) = \begin{cases} 0 & \text{falls } x = y, \\ 1 & \text{falls } x \neq y. \end{cases}$$

- 4. Why do you need the μ -recursion in your construction?
- 5. Is you construction in Kleene's normal form? If it is not, describe an (informal) procedure how one can turn it into Kleene's normal form.

Problem 23. Let P be the following program.

END;

Similar to the construction in the lecture notes, let $f_P : \mathbb{N}^2 \to \mathbb{N}^2$ be the function that maps the given $(0, x_1)$ at the start of P to the values (x_0, x_1) after the execution of the program P. Show that f_P is primitive recursive by translating the loop program into a primitive recursive definition for f_P . Follow the steps given in the lecture notes.

Compute $f_P(0,1)$ via your primitive recursive definition and compare it with the result you get from executing P with input $x_1 = 1$.

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Problem 24. Let $q: \mathbb{N}^2 \to \mathbb{N}, (x, y) \mapsto x \cdot x$ (sic!) and $u: \mathbb{N}^2 \to \mathbb{N}$,

$$u(x,y) = \begin{cases} 0 & \text{if } x = y, \\ 1 & \text{if } x \neq y, \end{cases}$$

be given primitive recursive functions. Let $r: \mathbb{N}^2 \to \mathbb{N}$ be defined by

$$r(x) = (\mu p)(x)$$
 minimization
 $p(y, x) = u(q(y, x), \operatorname{proj}_2^2(y, x))$ composition

Informally we have

$$r(x) = \min_{y} \{ y \in \mathbb{N} \, | \, u(q(y, x), x)) = 0 \, \}$$

Similar to the treatise in the lecture notes, construct a while program that computes r. For simplicity, you are allowed to write statements such as $x_k = q(x_i, x_j)$ and $x_k = u(x_i, x_j)$ into your program. What will your program compute if it is started with input $x_1 = 2$?

Problem 25. Let f be defined as

$$f(n) = \begin{cases} 3n+1 & \text{if } n \text{ is odd,} \\ \frac{n}{2} & \text{if } n \text{ is even.} \end{cases}$$

Now one can imagine the following process. Choose a positive natural number and apply f to it. If the result is 1 then terminate the process, otherwise apply f again and iterate the process until the result is 1.

Let ν be the function that takes a positive natural number x as input and returns the number of iterations in the above process until its termination, i.e. if the process terminates then

$$f^{\nu(x)}(x) = \underbrace{f(f(\cdots f(x)) \cdots)}_{\nu(x) \text{-fold}} = 1.$$

Show that ν is a recursive function. You may use any theorems from the lecture notes and you can use the (primitive recursive) function pred : $\mathbb{N} \to \mathbb{N}$,

pred(x) =
$$\begin{cases} 0 & \text{if } x = 0, \\ x - 1 & \text{if } x > 0. \end{cases}$$

Bonus task (very difficult): Is ν primitiv recursive?

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