## Problems Solved:

| 16 | 17 | 18 | 19 | 20 |
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## Name:

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## Problem 16.

Definition 1 (RAM computable). We say that a partial function $f: \mathbb{N} \rightarrow_{P} \mathbb{N}$ is $R A M$ computable if there exists a RAM $R$ that such that

- $R$ terminates for input $n \in \mathbb{N}$ if and only if $n \in \operatorname{domain}(f)$;
- $R$ terminates for input $n \in \mathbb{N}$ with output $n^{\prime}$ if and only if $n^{\prime}=f(n)$.

Show that every loop computable function is also RAM computable by describing how the loop program computing the function can be translated to a RAM program.
Problem 17. Give reasons for your answers.

1. Let $R$ be a RAM that reads exactly one number from its input tape and always terminates with 0 or 1 written on its output tape. Is there a function $f: \mathbb{Z} \rightarrow \mathbb{Z}$ such that $f(x)=y$ if and only if the input was $x$ and after termination $y$ is on the output tape of $R$ ?
2. Let $f: \mathbb{Z} \rightarrow \mathbb{Z}$ be a function. Is there always a RAM $R$ such that $R$ terminates on every input and that $R$ with input $z \in \mathbb{Z}$ has written $f(z)$ to its output tape?

Problem 18. Provide a loop program that computes the function $f(n)=$ $\sum_{k=1}^{n} k(k+1)$, and thus show that $f$ is loop computable.
Problem 19. Write down explicit loop programs for $s$ and $d$.

1. Show by using only the Definition of a loop program (Def. 23 in the lecture notes, Section 3.2.2) that the function

$$
s\left(x_{1}, x_{2}\right)= \begin{cases}1 & \text { if } x_{1}<x_{2} \\ 0 & \text { otherwise }\end{cases}
$$

is loop computable. I.e. give an explicit loop program for $s$.
Note that it is not allowed to use an abbreviation like
$\mathrm{x}_{\mathrm{i}}:=\mathrm{x}_{\mathrm{j}}-\mathrm{x}_{\mathrm{k}} ;$
2. Write a loop program that computes the function $d: \mathbb{N} \rightarrow \mathbb{N}$ where $d\left(x_{1}, x_{2}\right)$ is $k \in \mathbb{N}$ such that $k \cdot\left(x_{2}+1\right)=x_{1}+1$ if such a $k$ exists. The result is $d\left(x_{1}, x_{2}\right)=0$, if a $k$ with the above property does not exist.
For simplicity in the program for $d$, you are allowed to use a construction like the following (with the obvious semantics) where $P$ is an arbitrary loop program.

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IF }\mp@subsup{\textrm{x}}{\textrm{i}}{}<\mp@subsup{\textrm{x}}{\textrm{j}}{}\mathrm{ THEN P END;
```

Problem 20. Suppose $P$ is a while-program that does not contain any WHILE statements, but might modify the values of the variables $x_{1}$ and $x_{2}$.
Transform the following program into Kleene's normal form.
Hint: first translate the program into a goto program, replace the GOTOs by assignments to a control variable, and add the WHILE wrapper.
$\mathrm{x}_{0} \quad:=0$
WHILE $\mathrm{x}_{1}$ DO
$\mathrm{x}_{1}:=\mathrm{x}_{1}-1$;
$\mathrm{x}_{2}:=\mathrm{x}_{1}$;
WHILE $\mathrm{x}_{2}$ DO
P;
END;
END;
$\mathrm{x}_{0}:=\mathrm{x}_{0}+1$

