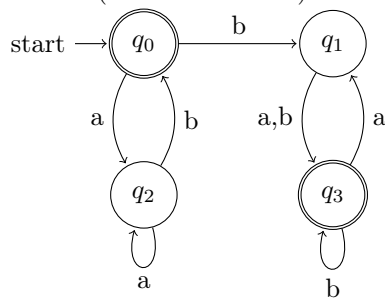


**Problems Solved:**

11	12	13	14	15
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**Name:****Matrikel-Nr.:**

**Problem 11.** Let  $M_1$  be the DFMSM with states  $\{q_1, q_2, q_3, q_4\}$  whose transition graph is given below. Determine a regular expression  $r$  such that  $L(r) = L(M_1)$ . Show the *derivation* of the the final result by the technique based on Arden's Lemma (see lecture notes).



**Problem 12.** Let  $r$  be the following regular expression.

$$a \cdot a \cdot (b \cdot a)^* \cdot b \cdot b^*$$

Construct a nondeterministic finite state machine  $N$  such that  $L(N) = L(r)$ . Show the derivation of the result by following the technique presented in the proof of the theorem *Equivalence of Regular Expressions and Automata* (see lecture notes).

**Problem 13.** Let  $L$  be the language of properly nested strings over the alphabet  $\Sigma = \{[, ], \circ\}$ . A word  $w$  is *properly nested* if it contains as many opening as closing brackets and every prefix of  $w$  contains at least as many opening brackets [ as closing ]. (Example:  $\circ\circ[[]\circ]$  is properly nested, but  $\circ\circ[[]$  is not.) Show by means of the Pumping Lemma that  $L$  is not regular.

**Problem 14.** Write down explicitly a Turing machine  $M$  over  $\Sigma = \{0\}$  which computes the function  $d: \mathbb{N} \rightarrow \mathbb{N}$  given by  $d(n) = 2n$ .

Use unary representation: A number  $n$  is represented by the string  $0^n$  consisting of  $n$  copies of the symbol 0.

**Problem 15.** Write down explicitly an enumerator  $G$  such that  $\text{Gen}(G) = \{0^{2n} \mid n \in \mathbb{N}\}$ .

Since in the lecture notes it has not been *formally* defined, how a Turing machine with two tapes works, you may describe the transition function as

$$\delta: Q \times \Gamma \rightarrow Q \times \Gamma \times \{R, L\} \times (\Gamma \cup \{\boxtimes\})$$

in the following way: If  $G$  is in state  $q$  and reads the symbol  $c$  from the working tape, and

$$\delta(q, c) = (q', c', d, c'')$$

then  $G$  goes to state  $q'$ , replaces  $c$  by  $c'$  on the working tape and moves the working tape head in direction  $d$ . Moreover, unless  $c'' = \boxtimes$ , the symbol  $c''$  is

written on the output tape and the output tape head is moved one position forward. If, however,  $c'' = \boxtimes$ , nothing is written on the output tape and the output tape head rests in place.

Hint: There exists a solution with only 4 states.