Formal Semantics of Programming Languages Exercise 2 (May 13)

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The exercise is to be submitted by the deadline stated above as a report with a decent cover page (title of the course, your name, Matrikelnummer, email address) in one of the following forms:

- 1. either as a single PDF file uploaded in Moodle (no emails, please), or
- 2. as a stapled paper report handed out to me (in class or in my mailbox).

1 Expressions with Side effects

Take the following language of commands *C*, expressions *E*, numerals *N*, and identifiers *I*:

$$C ::= I := E | C_1; C_2 | \text{ if } (E_1 = E_2) C$$

$$E ::= I | N | E_1 + E_2 | \text{ exec } C \text{ result } E$$

The exec expression executes C and then returns the result of the evaluation of E. Correspondingly, the evaluation of an expression may alter the store.

- 1. Define a denotational semantics for this language.
- 2. Define a big-step operational semantics for this language with transitions $\langle E, s \rangle \rightarrow \langle n, s' \rangle$ (expression *E* evaluated in store *s* yields number *n* and store *s'*).
- 3. Formulate the statement "the operational semantics of E is sound with respect to the denotational semantics" and prove this statement.
- 4. Define a small-step operational semantics for this language.
- 5. Formulate the statement "the small-step operational semantics of E corresponds to its big-step semantics".